

# Finding collisions for SHA-1

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# The near-anniversary of not a birthday search

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- ▶ On 2017-01-15, the first (public?) SHA-1 collision was found
- ▶ ... Coming after the first *freestart* collision in Oct. 2015
- ▶ ... Coming after the first “theoretical” attack in 2005
- ▶ ... Coming after the first standardization of SHA-1 in 1995

Aim of this talk:

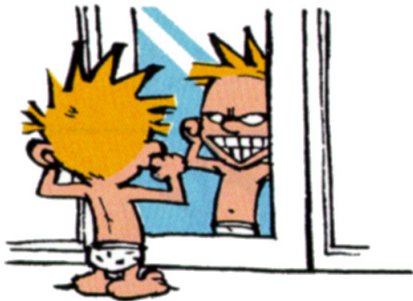
- ▶ What's a SHA-1 collision like? How do you **compute one**?
- ▶ How do you measure the “complexity” of such an attack?

# A simple collision

$h_0$	4e a9 62 69 7c 87 6e 26 74 d1 07 f0 fe c6 79 84 14 f5 bf 45
$M_1$	<u>7f</u> 46 dc <u>93</u> <u>a6</u> b6 7e <u>01</u> <u>3b</u> 02 9a <u>aa</u> <u>1d</u> b2 56 <u>0b</u> <u>45</u> ca 67 <u>d6</u> <u>88</u> c7 f8 <u>4b</u> <u>8c</u> 4c 79 <u>1f</u> <u>e0</u> 2b 3d <u>f6</u> <u>14</u> f8 6d <u>b1</u> <u>69</u> 09 01 <u>c5</u> <u>6b</u> 45 c1 <u>53</u> <u>0a</u> fe df <u>b7</u> <u>60</u> 38 e9 <u>72</u> <u>72</u> 2f e7 <u>ad</u> 72 8f 0e <u>49</u> <u>04</u> e0 46 <u>c2</u>
$h_1$	8d 64 <u>d6</u> <u>17</u> ff ed <u>53</u> <u>52</u> eb c8 59 15 5e c7 eb <u>34</u> <u>f3</u> 8a 5a 7b
$M_2$	<u>30</u> 57 0f <u>e9</u> <u>d4</u> 13 98 <u>ab</u> <u>e1</u> 2e f5 <u>bc</u> <u>94</u> 2b e3 <u>35</u> <u>42</u> a4 80 <u>2d</u> <u>98</u> b5 d7 <u>0f</u> <u>2a</u> 33 2e <u>c3</u> <u>7f</u> ac 35 <u>14</u> <u>e7</u> 4d dc <u>0f</u> <u>2c</u> c1 a8 <u>74</u> <u>cd</u> 0c 78 <u>30</u> <u>5a</u> 21 56 <u>64</u> <u>61</u> 30 97 <u>89</u> <u>60</u> 6b d0 <u>bf</u> 3f 98 cd <u>a8</u> <u>04</u> 46 29 <u>a1</u>
$h_2$	1e ac b2 5e d5 97 0d 10 f1 73 69 63 57 71 bc 3a 17 b4 8a c5
$h_0$	4e a9 62 69 7c 87 6e 26 74 d1 07 f0 fe c6 79 84 14 f5 bf 45
$M_1 \oplus \Delta_1$	<u>73</u> 46 dc <u>91</u> <u>66</u> b6 7e <u>11</u> <u>8f</u> 02 9a <u>b6</u> <u>21</u> b2 56 <u>0f</u> <u>f9</u> ca 67 <u>cc</u> <u>a8</u> c7 f8 <u>5b</u> <u>a8</u> 4c 79 <u>03</u> <u>0c</u> 2b 3d <u>e2</u> <u>18</u> f8 6d <u>b3</u> <u>a9</u> 09 01 <u>d5</u> <u>df</u> 45 c1 <u>4f</u> <u>26</u> fe df <u>b3</u> <u>dc</u> 38 e9 <u>6a</u> <u>c2</u> 2f e7 <u>bd</u> 72 8f 0e <u>45</u> <u>bc</u> e0 46 <u>d2</u>
$h_1$	8d 64 <u>c8</u> <u>21</u> ff ed <u>52</u> <u>e2</u> eb c8 59 15 5e c7 eb <u>36</u> <u>73</u> 8a 5a 7b
$M_2 \oplus \Delta_2$	<u>3c</u> 57 0f <u>eb</u> <u>14</u> 13 98 <u>bb</u> <u>55</u> 2e f5 <u>a0</u> <u>a8</u> 2b e3 <u>31</u> <u>fe</u> a4 80 <u>37</u> <u>b8</u> b5 d7 <u>1f</u> <u>0e</u> 33 2e <u>df</u> <u>93</u> ac 35 <u>00</u> <u>eb</u> 4d dc <u>0d</u> <u>ec</u> c1 a8 <u>64</u> <u>79</u> 0c 78 <u>2c</u> <u>76</u> 21 56 <u>60</u> <u>dd</u> 30 97 <u>91</u> <u>d0</u> 6b d0 <u>af</u> 3f 98 cd <u>a4</u> <u>bc</u> 46 29 <u>b1</u>
$h_2$	1e ac b2 5e d5 97 0d 10 f1 73 69 63 57 71 bc 3a 17 b4 8a c5

## A comic application

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```
>sha1sum *.pdf
```

```
23aa25d9e0449e507a8b4c185fdc86c35bf609bc calvin.pdf
```

```
23aa25d9e0449e507a8b4c185fdc86c35bf609bc hobbes.pdf
```

## SHA-1 collisions recap

On the way to full practical attacks

What complexity for an attack

Conclusion & Future work

# SHA-1 quick history

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## Secure Hash Standard “SHA-1”

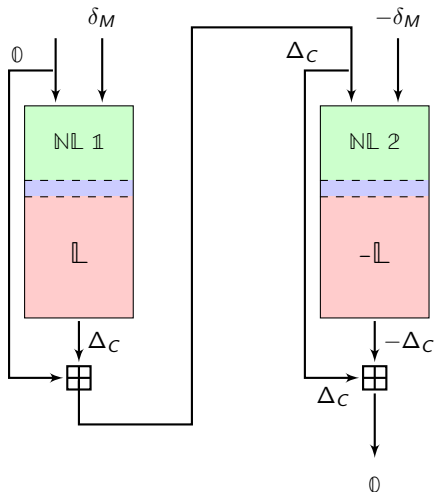
- ▶ Standardized by NIST in Apr. 1995
- ▶ Similar to MD4/5
  - ▶ Merkle-Damgård domain extender
  - ▶ Compression function = ad hoc block cipher in Davies-Meyer mode
  - ▶ Unbalanced Feistel network, 80 steps
- ▶ Quick fix of “SHA-0” (May 1993)
- ▶ Hash size is 160 bits  $\Rightarrow$  collision security should be 80 bits

That's nice, but we want to attack it!

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# A two-block attack in a picture





# The result

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- ▶ SHA-1 is **not collision-resistant** (Wang, Yin & Yu, 2005)
- ▶ Attack complexity  $\equiv 2^{69}$  (theoretical)
- ▶ Eventually improved to  $\equiv 2^{61}$  (ditto, Stevens, 2013)

# The attack process

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- 1 Pick a **linear path**
- 2 Find a **non-linear path** (first block)
- 3 Find **accelerating techniques** (first block)
- 4 Compute a *near-collision* (a solution for  $(0, \delta_M) \rightarrow \Delta_C$ )
  - ▶ Possible expected wall time estimation (first block)
- 5 Find a non-linear path (second block)
- 6 Find accelerating techniques (second block)
- 7 Compute a *collision* (a solution for  $(\Delta_C, -\delta_M) \rightarrow -\Delta_C$ )
  - ▶ Possible expected wall time estimation (full attack)

# Wall time estimation

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Simple approach:

- ▶ Implement the attack
- ▶ **Measure** production rate  $\#A_{xx}/s$
- ▶ Multiply by probability that a solution  $A_{xx}$  extends to  $A_{80}$

Early variant (crude):

- ▶ Partial solutions for the differential path up to  $A_{16}$  are free
- ▶ For  $A_{17...??}$ , count *path conditions* v. accelerating technique “efficiency”
- ▶ Estimate the “critical” step  $A_{xx}$  & corresp. production rate
- ▶ Multiply by probability that a solution  $A_{xx}$  extends to  $A_{80}$

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## Best practical attack progress (2005-2011)

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- ▶ 2005 (Biham & al.): 40 steps (cost: “within seconds”)
- ▶ 2005 (Wang & al.): 58 steps (cost:  $\approx 2^{33}$  SHA-1 computations)
- ▶ 2006 (De Cannière & Rechberger): 64 (cost:  $\approx 2^{35}$ )
- ▶ 2007 (Rechberger & al.): 70 (cost:  $\approx 2^{44}$ )
- ▶ 2007 (Joux & Peyrin): 70 (cost:  $\approx 2^{39}$ )
- ▶ 2010 (Grechnikov): 73 (cost:  $\approx 2^{50.7}$ )
- ▶ 2011 (Grechnikov & Adinets): 75 (cost:  $\approx 2^{57.7}$ )

# 2014: time to improve things again!

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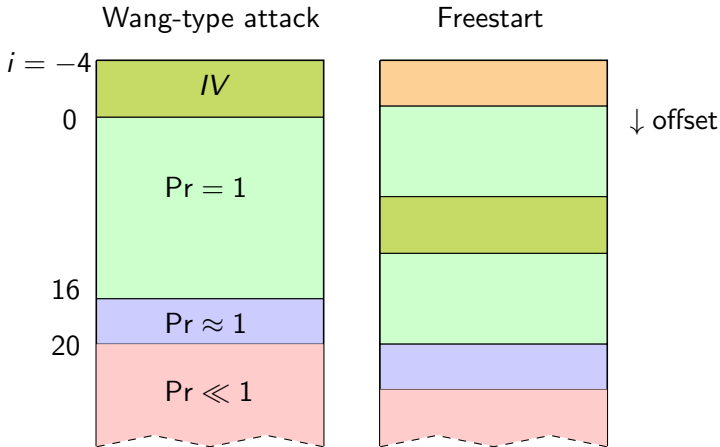
- ▶ Eventual objective: full practical collision??
- ▶ Significant intermediate step: full practical *freestart* collision?
  - ▶ Easier in principle, but is it the case?

⇒

- ▶ Search for a 76-step freestart collision (lowest # unattacked steps)
- ▶ Use the opportunity to develop a GPU framework

# The point of freestart (in a picture)

Internal state of SHA-1 ( $A_i$ )



# First results

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In Dec. 2014: a first **76-step freestart collision** (with Peyrin & Stevens)

- ▶ Right on time for the ASIACRYPT rump session :P
- ▶ Cost:  $\approx 2^{50}$  SHA-1 computations *on a GTX-970*  $\Rightarrow$  Freestart helps!
- ▶  $\Rightarrow$  **About 4 days on a single GPU** (what we did)
- ▶  $\Rightarrow$  About 1 day on a S\$ 3000 4-GPU machine



# Now what?

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WHAT IF WE TRIED  
MORE POWER?



# Objective: full compression function collision

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- ▶ Early (optimistic?) estimates: full freestart  $\approx 32\times$  more expensive than 76-step
- ▶ (Hard to know for sure w/o implementing it)
- ▶  $\Rightarrow$  buy (a bit) **more GPUs!**
- ▶ + **develop a new attack** (“sadly” necessary)
  - ▶ Update path search tools
  - ▶ Settle on a linear path
  - ▶ Generate new attack parameters
  - ▶ Program the attack again
  - ▶ ...

# Let's do this!

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Figure: Part of a homemade cluster to be

## Second results

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In Sep. 2015: a first 80-step (full) **freestart collision** (with Stevens & Peyrin)

- ▶ Right on time for EUROCRYPT submissions :P
- ▶ cost:  $\approx 2^{57.5}$  SHA-1 computations on a GTX-970
  - ▶ A bit more than expected
- ▶  $\Rightarrow$  **About 680 days on a single GPU**
- ▶ ... or 10 days on a 64-GPU cluster (what we did)
- ▶ ... or US\$ 2000 of the cheapest Amazon EC2 instances

## Some early impact

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- ▶ SHA-1 TLS certificates are *not extended* through 2016 by CA/Browser forum actors
  - ▶ Ballot 152 (Oct. 2015!) of the CA/Browser forum is withdrawn
- ▶ Some major browsers (Edge, Firefox) sped-up deprecation/security warnings
- ▶ But (some) continued use in Git, company-specific certificates (e.g. Facebook until Dec. 2016, Cloudflare), etc.
  - ▶ Mostly because of legacy issues

# Now what?

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WHAT IF WE TRIED  
MORE POWER?



## Objective: full hash function collision

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- ▶ Early (optimistic?) estimates: full collision  $\approx 50\times$  more expensive than full freestart
- ▶ (Hard to know for sure w/o implementing it)
- ▶  $\Rightarrow$  buy a lot more GPUs? (No)
- ▶  $\Rightarrow$  get help from GPU-rich people/companies? (Yes)
- ▶ + develop a new attack
- ▶ + add some cool exploitation features!

# Let's do this!

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## A CWI/Google collaboration

- 1 Prepare a prefix for future colliding PDFs
- 2 Compute a first (actually two) near-collision block(s)
  - ▶ Done on CPU
- 3 Compute a second near-collision  $\Rightarrow$  the final one!!
  - ▶ Done on GPU
- 4 Profit! Enjoy!
  - ▶ cost:  $\approx 2^{63}$  SHA-1 computations
    - ▶ A bit more/less than expected
  - ▶  $\Rightarrow$  about 6 500 CPU-year + 100 GPU-year
  - ▶ ... or US\$ 100K+ of the cheapest Amazon instances (second block only)



## Some more impact

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- ▶ Finally got Git planning to move away from SHA-1
- ▶ Unwittingly broke SVN for a time
- ▶ Further deprecation of SHA-1 certificates

SHA-1 collisions recap

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# Absolute cost v. “complexity”

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- ▶ Determining the complexity of generic attacks is “easy”
- ▶ E.g.  $\Theta(2^{n/2})$  for collisions on  $n$ -bit hash functions
  - ▶ Efficiently parallelizable (van Oorschot & Wiener, 1999)
- ▶ What about dedicated attacks?
  - ▶ Implement and measure?

A typical metric for cryptanalysis complexity:

- 1 Estimate the cost of an attack *on some platform*
- 2 Divide by the cost of computing the attacked function
- 3 Voilà

# A '76 complexity example

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Example: 76-step freestart collision

On a GTX-970:

- ▶ Expected time to collision = 4.4 days
  - ▶ 0.017 solution up to  $A_{56}/s$
- ▶  $\approx 2^{31.8}$  SHA-1 compression function/s
- ▶  $\Rightarrow 4.4 \times 86400 \times 2^{31.8} \approx 2^{50.3}$

BUT on an Haswell Core i5:

- ▶ Expected time to collision = 606 core days
  - ▶ 0.000124 solution up to  $A_{56}/s$
- ▶  $\approx 2^{23.5}$  SHA-1 compression function/s
- ▶  $\Rightarrow 606 \times 86400 \times 2^{23.5} \approx 2^{49.1}$
- ▶ Yet **much slower & less energy efficient!!**

## A full hash example

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Complexity for the full hash function (second block) collision:

- ▶  $2^{62.1}$  on K80, or
- ▶  $2^{62.8}$  on K20/40, or
- ▶  $2^{63.4}$  on GTX-970

Further code tuning/optimization may again change figures!

## Some more issues

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- ▶ Variation between CPU/GPU and optimized/unoptimized is not so large
  - ▶ About  $\times 2-4$
- ▶ What about reconfigurable/dedicated hardware?
  - ▶ FPGA/ASICs are fast and energy efficient
  - ▶  $\Rightarrow$  Well-suited to generic attacks!
  - ▶ But what about complex ones???
- ▶ No reason for a generic attacker to use CPU/GPU over FPGA/ASIC
  - ▶ Potential increased development cost well worth it!
- ▶ What does a dedicated attack really improve on??

# GPU v. ASIC brute force estimates

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One generic SHA-1 collision in one year  $\approx 2^{80}$  hash computations  
On GPU:

- ▶  $\approx 12.6$  million GPUs @  $2^{31.5}$  hashes/s
- ▶  $\approx 3.1$  GW 'round the clock (just the GPUs @ 250 W each)
  - ▶ A couple of dedicated nuclear powerplant needed

On ASIC (estimates courtesy of BTC mining hardware)

- ▶  $\approx 2900$  devices @  $2^{43.6}$  hashes/s (Antminer S9-like)
- ▶  $\approx 4$  MW 'round the clock (at 1400 W each)
  - ▶ About a large wind turbine needed (with the wind)

## An alternative cost measure: The fun calorie

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- ▶ Introduced by A. Lenstra, Kleinjung & Thomé (2013):  
How much energy is wasted needed by an attack?
- ▶ Energy unit: “fun calorie”  
What volume of standard water can you boil (instead)?
- ▶ Used to estimate e.g. RSA-768 security  
⇒ 2 olympic pool security (Kleinjung et al., 2010)



## Some complexity figures

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SHA-0 collision (MP08)	$\approx$ <b>teaspoon sec.</b> ( $2.5 \times 10^{-3}L$ )
SHA-1 76' fs.	$\approx$ <b>4 shower sec.</b> (320L)
SHA-1 fs.	$\approx$ <b>580 shower sec.</b> ( $4.5 \times 10^4L$ )
SHA-1 2 <sup>nd</sup> block (ded, GPU)	$\approx$ <b>1 pool sec.</b> ( $2.5 \times 10^6L$ )
RSA-768 (K+10)	$\approx$ <b>2 pool sec.</b> ( $5 \times 10^6L$ )
SHA-1 1 <sup>st</sup> block (ded, CPU)	$\approx$ <b>3 pool sec.</b> ( $7.5 \times 10^6L$ )
DL-768 (K+17)	$\approx$ <b>6 pool sec.</b> ( $1.5 \times 10^7L$ )
SHA-0/1 (gen, ASIC) <sup>†</sup>	$\approx$ <b>0.004 rain sec.</b> <sup>‡</sup> ( $3.5 \times 10^8L$ )

(Ignoring CPU improvements between 2010 and today)

<sup>†</sup>: Estimate

<sup>‡</sup>: *dagelijkse neerslagverdampingenergiebehoeftezekerheid*

## In the end...

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- ▶ Full-GPU dedicated SHA-1 attack:  $\approx 1$  pool sec.
- ▶  $\Rightarrow \approx 100\times$  better than dedicated hardware (conjectured)
- ▶ Quite less than  $2^{80-63} \approx 130\,000$

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## Potential future work

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- ▶ Computing a *chosen-prefix* collision
  - ▶ More exploitation
- ▶ Computing a collision for the SHA-1||MD5 combiner
  - ▶ Wouldn't break SVN?
- ▶ Designing a SHA-1-based crypto-currency
  - ▶ Get shiny mining hardware!

## For more details

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- ▶ **The papers:** Eprints 2015/530, 2015/967, 2017/190
- ▶ **The attack code:** [https://github.com/cr-marcstevens/sha1\\_gpu\\_nearcollisionattacks](https://github.com/cr-marcstevens/sha1_gpu_nearcollisionattacks)
- ▶ **Marc's talk** @ CRYPTO'17
- ▶ **Ange's talk** @ BlackAlps'17

C'est fini!

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